

King Fahd University of Petroleum & Minerals

College of Computer Science and Engineering Information and Computer Science Department Second Semester 202 (2020/2021)

> ICS 202 – Data Structures Midterm Exam Tuesday, 16th March 2021 Time: 120 minutes

Name:			 		
ID#					
Section	:				

- 1. This exam consists of 13 pages including (1) title page, (2) statement to ensure non-cheating and (13) reference sheet
- 2. It is required to sign your name on the (2) statement to ensure non-cheating. You may either print and sign it or use an MS-Paint Signature (please keep it ready before the exam).
- 3. Please solve this exam on this template. Use either (a) a printed version, solve by hand and scan or (b) solve it on your PC, or (c) use a tablet-stylus combination.
- 4. Please upload your solution as a pdf.
- 5. There will be a small penalty for solving this exam on plain paper. No excuses in this regard.

Question #	Max Points	Points Earned	Comments
1 [Linked Lists]	20		
2 [Stacks, Queues]	20		
3 [Complexity]	10		
4 [Recursion]	30		
5 [Binary Trees]	20		
Total	100		

Approved "Student Declaration Statement for Exam Integrity"

We wish you all the best in the exam.

Please read carefully and accept the following before proceeding to the exam.

I declare that:

- I will complete this assessment entirely by myself without taking any assistance whatsoever from a person, resource or tool other than what is explicitly permitted within the regulations prescribed for this assessment.
- I will not disclose, share or discuss the material of this assessment in any form with anyone except what has been authorized.
- I will uphold the highest standards of honesty and integrity for completing this assessment to the best of my knowledge.
- I am fully accountable for all rules pertaining to this assessment.
- I understand that ignorance of a rule or a principle is not an excuse for any misconduct.
- I understand that engaging in an act of a misconduct would result in imposition of penalties such as failing this course and/or a dismissal from the University.

مع امنياتنا لكم بالتوفيق والنجاح في هذا الاختبار نرجو قراءة النص التالي بعناية والموافقة عليه قبل البدء بالاختبار أقرب:

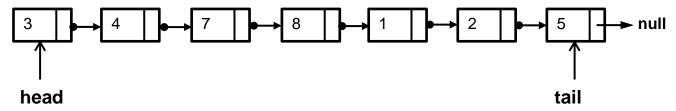
سرب. أن أنهي الاختبار بأكمله بنفسي وبدون أن أتلقى أي مساعدة من أي شخص أو من أي مصدر بخلاف ما هو منصوص عليه بوضوح في لوائح هذا الاختبار.

- أن لا أقوم بنشر أو بمشاركة أو بمناقشة أي من محتويات هذا الاختبار مع أي شخص باستثناء ما سمح لي فيه.
 - الالتزام بأعلى معايير الأمانة والنزاهة في إكمال هذا الاختبار.
 - أن أكون مسؤو لا مسؤولية تامة عن تطبيق جميع الأنظمة المتعلقة بهذا الاختبار.
 - أن جهلي بأي من أنظمة أو متطلبات هذا الاختبار لا تعفيني من المسؤولية تجاه أي تصرف غير مقبول.
- علمي بأن أي تصرف من قبلي يخل بالامانة او النزاهة في اخذ الاختبار قد يعرضني إلى العقوبات التي تقرها الجامعة والتي قد تصل إلى الرسوب في المقرر و/أو الفصل من الجامعة.

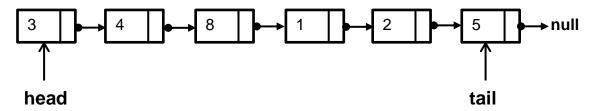
Name:	
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- Q. 1: (a) [14 points] Consider a singly linked list represented by the class SLL<T> as shown in the reference sheet. Design and implement the following two methods:
 - (i) deleteThird which deletes the third element of a singly linked list.
 - (ii) deleteThirdLast which deletes the third last of a singly linked list.

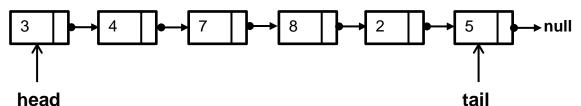
The effect of these methods is shown here (initial linked list):



[After applying myList.deleteThird()]



[After applying myList.deleteThirdLast()]



Do not use the methods addToHead, addToTail, deleteFromHead, deleteFromTail. Instead, directly manipulate list nodes/pointers for both the methods.

Make sure to take care of *all* special cases

(If the third or third last element does not exist, just return from the method/s).

```
(i) public void deleteThird() {
    if(head == null || head.next == null || head.next.next == null)
        return;
    if(head.next.next == tail) //only three elements
        tail = head.next; //tail becomes the second element
    head.next.next = head.next.next.next; //takes care of all cases
}
(ii) public void deleteThirdLast() {
```

<pre>if(head == null head.next == null head.next.next == null)</pre>
return;
<pre>if(head.next.next == tail) // only three elements</pre>
<pre>{ head = head.next; return; } //delete the first element</pre>
SLLNode <t> prev = head;</t>
<pre>while(prev.next.next != tail) prev = prev.next;</pre>
<pre>prev.next = prev.next.next;</pre>
return;

Q. 1(b) [6 points] What is the big-O time complexity of both your methods in terms of list size n.

Big-O Complexity of deleteThird(): O(1)
since no traversal of list is involved
Big-O Complexity of deleteThirdLast(): O(n)
since list is traversed to find the third last element

Q. 2 (a) [10 points] Write a method **public static boolean isPalindrome(String s)** that determines whether an input string **s** is a palindrome or not. [A palindrome is a string that reads the same forwards and backwards. For example: *level* is a palindrome, but *lever* is not]. Do not use arrays or any other data structure for this program except Stacks. Consider using multiple stacks.

<pre>public static boolean isPalindrome(String s) {</pre>
<pre>Stack s1 = new Stack(); String s2 = new String();</pre>
<pre>for(int ix = 0; ix < s.length(); ix++)</pre>
s1.push(s.charAt(ix) + "");
while(!s1.isEmpty())
s2 = s2 + s1.pop();
return s1.equals(s2);
}

Q. 2 (b) [10 points] Given the following infix expression:

(i) Give the equivalent postfix expression.

(ii) Using a stack, evaluate this postfix expression (Give contents of stack at each stage). [The first two rows are just examples for the expression 2 3 *]

Stack Contents	Operations
2	Push 2, Push 3, Remaining Expression: *
<u> 6 </u>	Pop (3), Pop (2), 3*2 = 6, Push(6)

Stack Contents	Operations
8 10	Push 10, Push 8, Remaining Expression: *
<u> 80 </u>	Pop (8), Pop (10), 10 * 8 = 80, Push(80)
4 8 <u> 80 </u>	Push 8, Push 4, Remaining Expression: /
2 80	Pop (4), Pop (8), 8 / 4 = 2, Push(2)
3 2 80	Push 3, Remaining Expression: —
-1 <u>80</u>	Pop (3), Pop (2), $2-3=-1$, Push(-1), Remaining Operation +
<u> 79 </u>	Pop (-1) , Pop (80) , $80 + -1 = 79$, Push (79)
	Pop (79): Final Answer

```
Q. 3: [7 + 3 = 10 \text{ points}].
```

- (a) Given the following method, how many times is MyStatement executed as a function of n.
- (b) Give the big-O complexity of this code fragment in terms of O(n).

```
for(i = 1; i <= \sqrt{n}; i++) { sum[i] = 0; for (j = 1; j <= i<sup>3</sup>; j++) sum[i] = sum[i] + j; // MyStatement } return true;
```

$\sum_{i=1}^{\sqrt{n}} \sum_{j=1}^{i^3} 1 = \sum_{i=1}^{\sqrt{n}} i^3 = [\sqrt{n} (\sqrt{n} + 1) / 2]^2 = \frac{n(n + 2\sqrt{n} + 1)}{4}$	
Big-O complexity: O(n²)	

Q. 4 [30 points: 2 + 15 + 13 = 30 points]

(a) Write the recurrence relation that represents the number of **additions T(n)** as a function of n in the following method:

```
public static int myMethod(int n){
    if(n == 0)
        return 0;
    else{
        System.out.println(n);
        return myMethod(n - 2) + n;
    }
}
```

Note: DO NOT EXPAND THE RECURRENCE RELATION

(a)

T(n) = 0 for n = 0, T(n) = T(n-2) + 1, otherwise for n > 0.

(b) [15 points] The running time **T(n)** of an algorithm is represented by the following recurrence relation:

$$T(0) = a$$

$$T(n) = T(n-1) + \frac{n}{2} + b$$

$$\forall n > 0$$

Where a and b are constants. Solve the recurrence relation by **iteration** and then determine the big-O complexity of the algorithm.

You may find the following summation formulae useful:

$\sum_{i=1}^{n} i = \frac{n(n+1)}{2}$	$\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$	$\sum_{i=0}^{k-1} \frac{1}{2^i} = 2 - \frac{1}{2^{k-1}}$	$\sum_{i=0}^{k-1} 2^i = 2^k - 1$
---------------------------------------	---	--	----------------------------------

Solution:

$$T(n) = T(n-1) + \frac{n}{2} + b$$

$$= [T(n-2) + \frac{(n-1)}{2} + b] + \frac{n}{2} + b$$

$$= T(n-2) + \frac{1}{2}((n-1) + n) + 2b$$

$$= [T(n-3) + \frac{(n-2)}{2} + b] + \frac{1}{2}((n-1) + n) + 2b$$

$$= T(n-3) + \frac{1}{2}((n-2) + (n-1) + n) + 3b$$

$$= T(n-k) + \frac{1}{2}((n-k+1) + ... + (n-2) + (n-1) + n) + kb$$

By substituting k = n,

$$T(n) = T(0) + \frac{1}{2}(1+2+...+(n-1)+n) + nb$$

$$= a + \frac{1}{2}\sum_{i=1}^{n} i + nb = a + \frac{1}{2}(\frac{n(n+1)}{2}) + nb = a + \frac{n(n+1)}{4} + nb = a + \frac{n^2}{4} + \frac{n}{4} + nb$$

Therefore, $T(n) = O(n^2)$

(c) Write a method: **public static getMax(int[] array)** that calls a private static recursive method **getMax** which returns the maximum value in the **array**.

```
public static int getMax(int[] array){
   int currentMax = array[array.length - 1];
   return getMax(array, currentMax, array.length - 1);
}
```

(d) Write the recursive method that is called by the method you wrote in (c)

```
private static int getMax(int[] array, int currentMax, int index){
    if(index == -1)
        return currentMax;
    else if(array[index] > currentMax)
        return getMax(array, array[index], index - 1);
    else
        return getMax(array, currentMax, index - 1);
}
```

Q. 5: (a) Given the following BSTNode class:

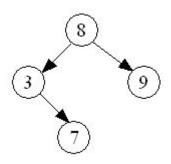
```
public class BSTNode<T extends Comparable<? super T>> {
    protected T el;
    protected BSTNode<T> left, right;
    public BSTNode() {
        left = right = null;
    }
    public BSTNode(T el) {
        this(el,null,null);
    }
    public BSTNode(T el,BSTNode<T> lt, BSTNode<T> rt){
        this.el = el; left = lt; right = rt;
    }
}
```

And the following instance methods of **BinarySearchTree** class:

```
public void myTraversal(){
    myTraversal(root);
}

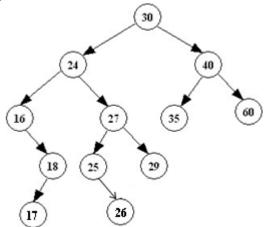
private void myTraversal(BSTNode node){
  if(node != null){
    System.out.print(node.el + " ");
    myTraversal(node.left);
    myTraversal(node.right);
    System.out.print(node.el + " ");
}
```

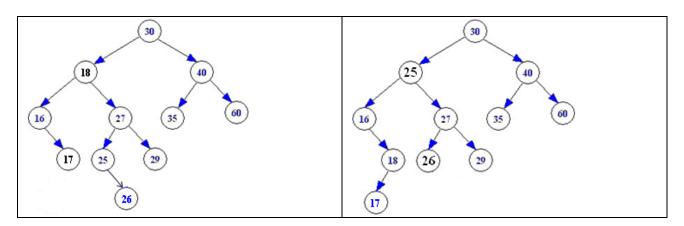
What is the output of: **tree.myTraversal()**; if **tree** is the following BinarySearchTree?



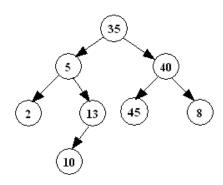
8 3 7 7 3 9 9 8

(b) Draw the resulting Binary search tree after deleting 24 by \underline{ANY} deletion by **copying method** from the following BST:





(c) Give the inorder, preorder, and postorder traversals of the following BinaryTree:



Traversal type	traversal
Inorder	2, 5, 10, 13, 35, 45, 40, 8
preorder	35, 5, 2, 13, 10, 40, 45, 8
postorder	2, 10, 13, 5, 45, 8, 40, 35

Quick Reference Sheet

```
public class SLLNode<T> {
   public T info;
   public SLLNode<T> next;
  public SLLNode();
  public SLLNode(T el)
  public SLLNode(T el, SLLNode<T> ptr);
public class SLL<T> {
   protected SLLNode<T> head, tail;
  public SLL();
  public boolean isEmpty();
  public void addToHead(T el);
  public void addToTail(T el);
  public T deleteFromHead();
  public T deleteFromTail();
  public void delete(T el);
 public void printAll();
 public boolean isInList(T el);
public class DLLNode<T> {
   public T info;
   public DLLNode<T> next, prev;
  public DLLNode();
  public DLLNode(T el);
  public DLLNode(T el, DLLNode<T> n,
                       DLLNode<T> p);
public class DLL<T> {
   private DLLNode<T> head, tail;
  public DLL();
  public boolean isEmpty();
  public void setToNull();
  public void addToHead(T el);
  public void addToTail(T el);
 public T deleteFromHead();
 public T deleteFromTail();
 public void delete(T el);
 public void printAll();
 public boolean isInList(T el);
```

```
public class Stack<T> {
     private ...; // array or linked list
   public Stack();
   public Stack(int n);
   public void clear();
   public boolean isEmpty();
   public T topEl();
   public T pop();
   public void push(T el);
   public String toString();
public class Queue<T> {
     private ...; // array or linked list
  public Queue();
  public void clear();
  public boolean isEmpty();
  public T firstEl();
  public T dequeue();
  public void enqueue(T el);
  public String toString();
```

$$\sum_{i=1}^{n} i = \frac{n(n+1)}{2}$$

$$\sum_{i=1}^{n} i^{2} = \frac{n(n+1)(2n+1)}{6}$$

$$\sum_{i=1}^{n} i^{3} = \left(\frac{n(n+1)}{2}\right)^{2}$$

$$\sum_{i=1}^{n} x^{i} = \frac{x^{n+1} - 1}{n-1}$$